

# Isomorphisms and Normality of Cayley Digraphs of $A_5$

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## Abstract

It is proved that each 2-element generating set of  $A_5$  is a *CI*-subset and that the corresponding Cayley digraph is normal. It is furthermore proved that for each 3-element generating set of  $A_5$  the corresponding Cayley digraph is normal.

## 1 Introduction

Let  $G$  be a finite group and  $S$  a subset of  $G$  not containing the identity element 1. We define the *Cayley digraph*  $X = \text{Cay}(G, S)$  of  $G$  with respect to  $S$  by

$$\begin{aligned}V(X) &= G, \\E(X) &= \{(g, sg) \mid g \in G, s \in S\}.\end{aligned}$$

If  $S = S^{-1}$ , then the adjacency relation is symmetric and  $\text{Cay}(G, S)$  is called the *undirected* Cayley graph of  $G$  with respect to  $S$ . The group  $G$  acting by right multiplication (that is,  $g_R : x \mapsto xg$ ) is a subgroup of automorphisms of  $\text{Cay}(G, S)$  and acts transitively on vertices. We call  $G_R = \{g_R \mid g \in G\}$  the *right regular representation* of  $G$ . Let  $\text{Aut}(G, S) = \{\alpha \in \text{Aut}(G) \mid S^\alpha = S\}$ . Obviously  $\text{Aut}(X) \geq G_R \text{Aut}(G, S)$ . If  $G_R = \text{Aut}(X)$ , then  $X$  is called a digraphical regular representation (*DRR*) of  $G$  and a *DRR* of a group  $G$  is a normal Cayley graph of  $G$ .

Let  $A = \text{Aut}(X)$ . We have

**Lemma 1.1** ([2, Proposition 1.3])

- (1)  $N_A(G_R) = G_R \text{Aut}(G, S)$ ;
- (2)  $A = G_R \text{Aut}(G, S)$  is equivalent to  $G_R \triangleleft A$ .

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**Definition 1.2** *The Cayley (di)graph  $X = \text{Cay}(G, S)$  is called normal if  $G_R$ , the right regular representation of  $G$ , is a normal subgroup of  $\text{Aut}(X)$ .*

So, normal Cayley digraphs are just those which have the smallest possible full automorphism group. The following obvious result is a direct consequence of the above definition and Lemma 1.1.

**Lemma 1.3** *Let  $X = \text{Cay}(G, S)$  be the Cayley digraph of  $G$  with respect to  $S$ , and let  $A = \text{Aut}(X)$ . Let  $A_1$  be the stabilizer of the identity element 1 in  $A$ . Then  $X$  is normal if and only if every element of  $A_1$  is an automorphism of the group  $G$ .*

Let  $X = \text{Cay}(G, S)$  be the Cayley digraph of  $G$  with respect to  $S$ . Let  $\alpha \in \text{Aut}(G)$ . Then it is easy to see that  $\alpha$  is a graph isomorphism from  $\text{Cay}(G, S)$  to  $\text{Cay}(G, S^\alpha)$ . We call this kind of isomorphism between Cayley digraphs of  $G$  a trivial automorphism. The subset  $S$  is said to be a *CI-subset* of  $G$ , if for any graph isomorphism  $\text{Cay}(G, S) \cong \text{Cay}(G, T)$ , there exists an  $\alpha \in \text{Aut}(G)$  such that  $S^\alpha = T$ . In other words, that  $S$  is *CI* means that there are only trivial isomorphisms between  $\text{Cay}(G, S)$  and other Cayley digraphs of  $G$ .

The motivation for this paper comes from a survey of Xu [2] and an unpublished result of Li [7] which states that some 2-element generating sets of  $A_5$  are *CI*-subsets, and the corresponding Cayley digraph is normal. In order to make this paper self-contained, we prove Li's result in section 2 while in section 3 we give a further extension. The main results of this paper are the following two theorems.

**Theorem 1.4** (See [2]) *Each 2-element generating set of  $A_5$  is a CI-subset and the corresponding Cayley digraph is normal.*

This result was originally proved by X. Li. However, the proof of Theorem 1.4 in section 2 is independent of Li's.

**Theorem 1.5** *Let  $G = A_5$  and  $S = \{a, b, c\}$  be a 3-element generating set of  $G$  not containing the identity 1. Then  $X = \text{Cay}(G, S)$  is a normal Cayley digraph.*

In this paper the symbol  $G$  will always denote the group  $A_5$  and 1 will denote its identity. For  $x \in G$  we let  $o(x)$  denote the order of  $x$ . The symbol  $X$  will always denote a simple graph. By  $V(X)$ ,  $E(X)$  and  $A(X) = A$  we denote the vertex set, the edge set and the automorphism group of  $X$ , respectively. By  $A_v(X) = A_v$  we denote the stabilizer of the vertex  $v \in V(X)$ . For every set  $T$ ,  $1_T$  denotes the identity permutation on  $T$ .

The group and graph-theoretic notation and terminology used here are generally standard, and the reader can refer to [3] and [6] when necessary.

## 2 The Proof of Theorem 1.4

**Lemma 2.1** *Let  $S = \{a, b\}$  and  $T = \{a', b'\}$  be two 2-element generating subsets of  $G = A_5$ . If  $X = \text{Cay}(G, S) \cong \text{Cay}(G, T) = X'$  and if  $\min \{o(a), o(b)\} \leq 3$ , then there exists  $\alpha \in \text{Aut}(G)$  such that  $S^\alpha = T$ .*

**Proof** Let  $\alpha$  be a graph isomorphism. Without loss of generality we may assume that  $1^\alpha = 1$ , since Cayley digraphs are vertex-transitive. Hence  $S^\alpha = T$ . By renaming the elements of  $T$  if necessary, we may also assume that

$$a^\alpha = a', \quad b^\alpha = b'.$$

We use induction on  $n$  to show that  $(x_1x_2 \cdots x_n)^\alpha = x'_1x'_2 \cdots x'_n$ , where  $x_i = a$  or  $b$  for  $i = 1, \dots, n$ . (This implies that  $\alpha \in \text{Aut}(G)$ , as required.) We distinguish two cases.

Case 1:  $\min\{o(a), o(b)\} = 2$ .

Without loss of generality, we may assume that  $o(a) = 2$ . Then  $o(b) \neq 2$  since  $\langle a, b \rangle \cong A_5$ . Thus  $(a, 1) \in E(X)$  and  $(b, 1) \notin E(X)$  by definition of the Cayley digraph  $X$  and so  $(a', 1) \in E(X')$ . By definition of the Cayley digraph  $X'$ ,  $(a', b'a') \in E(X')$  and  $(a', (a')^2) \in E(X')$ . If  $b'a' = 1$ , this would contradict  $\langle a', b' \rangle \cong A_5$ , so it follows that  $(a')^2 = 1$  and thus  $(ba)^\alpha = b'a'$  and  $o(b') \neq 2$ . Using the fact that the graphs have indegrees and outdegrees equal to 2, we have that  $(ab)^\alpha = a'b'$  and  $(b^2)^\alpha = (b')^2$ . So  $(x_1x_2)^\alpha = x'_1x'_2$ .

Suppose  $n > 2$ . Set  $x = x_3x_4 \cdots x_n$  and  $x' = x'_3x'_4 \cdots x'_n$ . From the inductive assumption we can suppose that  $x^\alpha = x'$ ,  $(ax)^\alpha = a'x'$  and  $(bx)^\alpha = b'x'$ . An argument similar to the one above shows that  $(x_1x_2x)^\alpha = x'_1x'_2x'$ . Thus we have that for any positive integer  $n$ ,  $(x_1x_2 \cdots x_n)^\alpha = x'_1x'_2 \cdots x'_n$ .

Case 2:  $\min\{o(a), o(b)\} = 3$ .

Let  $o(a) = 3$ . As  $1 \mapsto a \mapsto a^2 \mapsto 1$  is a directed circuit of length 3 in  $X$  we have that  $1 \mapsto a' \mapsto (a')^2 \mapsto (a')^3$  must be a directed circuit of length 3 in  $X'$  because  $a'b'a' \neq 1$ ,  $(b')^2a' \neq 1$  and  $b'(a')^2 \neq 1$  (otherwise  $G = \langle a', b' \rangle$  is abelian). Thus  $(a')^3 = 1$  and  $(a^2)^\alpha = (a')^2$ ,  $(ba)^\alpha = b'a'$ .

If  $o(b) = 3$ , similarly we have  $(b^2)^\alpha = (b')^2$  and  $(ab)^\alpha = a'b'$ . If  $o(b) \neq 3$  then we also have  $(b^2)^\alpha = (b')^2$  and  $(ab)^\alpha = a'b'$  because there is a unique directed circuit of length 3 through  $b$  and  $b'$  in  $X$  and  $X'$  respectively.

Thus for  $n = 1$  or  $2$ ,  $(x_1x_2 \cdots x_n)^\alpha = x'_1x'_2 \cdots x'_n$ . Using the same method as in (1) completes the inductive step.

It follows that  $\alpha \in \text{Aut}(G)$ . □

We now prove Theorem 1.4.

First we show that each 2-element generating subset  $S$  of  $G$  is  $CI$ .

In Lemma 2.1, we checked that all 2-element generating subsets of  $G$  are  $CI$ -subsets except the case  $o(a) = o(b) = 5$ . Thus to prove our statement it suffices to prove that given two 2-element generating subsets  $S = \{a, b \mid o(a) = o(b) = 5\}$  and  $T = \{a', b'\}$  such that  $X = \text{Cay}(G, S)$  and  $X' = \text{Cay}(G, T)$  are isomorphic, there exists an  $\alpha \in \text{Aut}(G)$  such that  $S^\alpha = T$ .

Suppose that  $\alpha : X = \text{Cay}(G, S) \rightarrow \text{Cay}(G, T)$  is a graph isomorphism. Since Cayley digraphs are vertex-transitive, without loss of generality we may assume that  $1^\alpha = 1$ . Hence  $S^\alpha = T$ . By renaming the elements of  $T$  if necessary, we may also assume that

$$a^\alpha = a', \quad b^\alpha = b'.$$

As  $1 \mapsto a \mapsto a^2 \mapsto a^3 \mapsto a^4 \mapsto 1$  is a directed circuit of length 5 in  $X$  we have that  $1 \mapsto a' \mapsto (a^2)' \mapsto (a^3)'\mapsto (a^4)'\mapsto (a^5)'$  must be a directed circuit of length 5 in  $X'$ .

We now distinguish three cases:

(i)  $(b')^i(a')^j = 1$ , for  $0 < i \leq 4$  and  $i + j = 5$ .

Then  $(b')^{-i} \in \langle a' \rangle$  and so  $(b') \in \langle a' \rangle$  which is a contradiction. So (i) cannot happen.

(ii)  $(b'a')^2a' = 1$  or  $(a'b')^2a' = 1$ .

If  $(b'a')^2a' = 1$  then  $a' \in \langle b'a' \rangle$  and thus  $b' \in \langle b'a' \rangle$  which contradicts  $A_5 \cong \langle a', b' \rangle \leq \langle a'b' \rangle$ . Similarly  $(a'b')^2a' \neq 1$  and so (ii) cannot happen.

(iii)  $a'(b')^2(a')^2 = 1$  or  $a'(b')^3a' = 1$ .

If  $a'(b')^2(a')^2 = 1$ , then it is easy to check  $(b')^2 \neq 1$  and thus  $(b')^2 \in \langle a' \rangle$ . So  $b' \in \langle a' \rangle$ , a contradiction. Similarly  $a'(b')^3a' \neq 1$ .

By (i), (ii), (iii), we have  $(a')^5 = (a^5)' = 1$  and thus  $1 \mapsto a' \mapsto (a')^2 \mapsto (a')^3 \mapsto (a')^4 \mapsto (a')^5$  is a directed circuit of length 5 in  $X'$ . So  $(a')^i = (a^i)'$  for  $i=1,2,3,4,5$  and thus  $b'a' = (ba)'$ . Similarly we have  $(b')^2 = (b^2)'$  and  $a'b' = (ab)'$ .

Thus for  $n = 1$  or  $2$ ,  $(x_1x_2 \cdots x_n)^\alpha = x'_1x'_2 \cdots x'_n$ . Using the same inductive method as in the proof of Lemma 2.1, we can show that  $\alpha \in \text{Aut}(G)$ .

Now we show that  $X = \text{Cay}(G, S)$  is a normal Cayley digraph of  $G$ .

Let  $S = \{a, b\}$  be a 2-element generating subset of  $G$  and  $A = \text{Aut}(X)$ . Then  $A = A_1G_R$ . For each  $\phi \in A_1$ ,

$$\phi : X = \text{Cay}(G, S) \longrightarrow X' = \text{Cay}(G, S^\phi)$$

is graph isomorphism. As shown above,  $S$  is  $CI$ , so  $\phi \in \text{Aut}(G)$  and this implies  $\phi \in \text{Aut}(G, S)$ . It follows that  $A_1 = \text{Aut}(G, S)$ . By Lemma 1.1,  $\text{Cay}(G, S)$  is normal as required.  $\square$

### 3 The Proof of Theorem 1.5

The proof is organized into twelve Lemmas.

**Lemma 3.1** *Let  $G = A_5$  and  $S = \{a, b, c\}$  be a 3-element generating subset of  $G$ . Set  $X = \text{Cay}(G, S)$  and  $A = \text{Aut}(X)$ . Then  $X = \text{Cay}(G, S)$  is normal if the following conditions hold:*

(1) *for each  $\phi \in A_1$ ,  $\phi|_S = 1_S$  implies  $\phi|_{S^2} = 1_{S^2}$ ,*

(2) *for each  $\phi \in A_1$ ,  $\phi^2|_S = 1_S$ .*

**Proof** Let  $H_S$  denote the subgroup of  $A$  which fixes 1,  $a$ ,  $b$ , and  $c$ . First we show that  $H_S$  is trivial. Note that since  $A$  is transitive on  $V(X) = G$ , condition (1) applies to every vertex  $v$ , that is, for any  $\phi \in A_v(X)$ ,  $\phi|_{Sv} = 1_{Sv}$  implies  $\phi|_{S^2v} = 1_{S^2v}$ . Now let  $\phi \in H_S$ , then  $\phi \in A$ , and  $\phi|_S = 1_S$  so  $\phi|_{S^2} = 1_{S^2}$ . Let  $x \in S$ , then  $x^\phi = x$ , so  $\phi \in A_x(X)$ . Also  $Sx \subseteq S^2$  so  $\phi|_{Sx} = 1|_{Sx}$ . Hence  $\phi|_{S^2x} = 1_{S^2x}$ . Since this holds for all  $x \in S$ , we have  $\phi|_{S^3} = 1_{S^3}$ . By induction,  $\phi|_{S^t} = 1_{S^t}$  holds for any positive

integer  $t$ . Since  $G = \langle S \rangle$  and  $G$  is finite, we have that if  $\phi_S = 1_S$  then  $\phi = 1_G$ . This implies that if  $\phi \in H_S$ , then  $\phi = 1_G$  and so  $H_S = 1$  as required.

Since condition (2) says that for each  $\phi \in A_1$ ,  $\phi^2 \in H_S$ , we have  $\phi^2 = 1_G$  for all  $\phi \in A_1$ , so  $A_1$  is 2-group. Since  $A_1 \subseteq S_3 H_S = S_3$ ,  $|A_1| \leq 2$ . As  $A = A_1 G_R$  we have  $G_R \triangleleft A$ . Thus  $X = \text{Cay}(G, S)$  is normal.  $\square$

Using this lemma we can analyze the normality of a Cayley digraph of  $A_5$  in terms of its generating set.

Let  $S$  be a generating subset of  $G$  of cardinality 3 and let  $l, m, n$  be integers  $\geq 2$ . We call  $S$  an  $(l, m, n)$ -generating set of  $G$  if  $a, b, c \in S$ ,  $a^l = b^m = c^n = 1$  and  $G = \langle a, b, c \rangle$ .

**Lemma 3.2** *Let  $S$  be a  $(2, 3, 5)$ -generating set of  $G$ . Then  $X = \text{Cay}(G, S)$  is a normal Cayley digraph. Moreover  $X$  is a DRR of  $G$ .*

**Proof** Now  $A_5$  is by definition the group given by the presentation  $\langle a, b, c \mid a^2 = b^3 = c^5 = 1, ab = c \rangle$  (or  $\langle a, b, c \mid a^2 = b^3 = c^5 = 1, ac = b \rangle$ ). Thus if  $S = \{a, b, c\}$  as above, then  $C_1 = (1, a)$ ,  $C_2 = (1, b, b^2)$  and  $C_3 = (1, c, c^2, c^3, c^4)$  are unique cycles of lengths 2, 3 and 5 at the point 1. So for each  $\phi \in A_1$ , we have  $C_1^\phi = C_1$ ,  $C_2^\phi = C_2$  and  $C_3^\phi = C_3$ . Since  $C_1$ ,  $C_2$  and  $C_3$  are directed dicircuits, therefore  $\phi$  fixes  $C_1$ ,  $C_2$  and  $C_3$  pointwise, and so  $b^\phi = b$ ,  $c^\phi = c$  and  $a^\phi = a$ . It follows that  $A_1 = A_a = A_b = A_c$ . Finally because  $a, b, c$  generate the group  $G$ , therefore  $A_1 = A_g$  for any  $g \in G$  and so  $A_1 = 1$  and  $A = G_R$  as required.  $\square$

In the remaining part of this section we shall discuss the case  $o(a) = o(b) \neq o(c)$ . If  $\phi$  is a non-trivial graph automorphism which fixes the point 1, it must have the form

$$\phi|_S = (a, b), \quad c^\phi = c.$$

This shows that  $\phi^2|_S = 1_S$ . Applying Lemma 3.1 in this case, to prove that  $\text{Cay}(G, S)$  is a normal Cayley digraph of  $G$ , we need only check that the condition (1) of Lemma 3.1 holds, that is, we focus on testing which digraphs meet the condition: for  $\phi \in A_1$ ,

$$\text{if } \phi|_S = 1_S, \text{ then } \phi|_{S^2} = 1_{S^2}. \quad (3.1)$$

**Lemma 3.3** *Let  $S$  be a  $(3, 5, 5)$ -generating set of  $G$  or a  $(3, 3, 5)$ -generating set of  $G$ . Then  $\text{Cay}(G, S)$  is a normal digraph.*

**Proof** Suppose that  $S = \{b, c_1, c_2\}$  and  $G = \langle b, c_1, c_2 \rangle$  or  $S = \{b_1, b_2, c\}$  and  $G = \langle b_1, b_2, c \rangle$ , where  $b, b_1, b_2$  are elements of order 3 and  $c_1, c_2, c$  are elements of order 5 in  $G$ . Let  $S_1 = \{b, c_1\}$  and  $S_2 = \{b_1, c\}$ . Since a subgroup of  $A_5$  which contains an element of order 3 and an element of order 5 is  $A_5$ , so  $G = \langle S_i \rangle$  for  $i = 1, 2$ . If  $\phi|_S = 1_S$  for  $\phi \in A_1$  then  $\phi$  induces an action on  $S_i$  for  $i = 1, 2$ . By Theorem 1.4  $\text{Cay}(G, S_i)$  is a normal Cayley digraph. Hence  $\phi \in \text{Aut}(G, S_i)$ , for  $i = 1, 2$ . In either case  $\phi$  is an automorphism of  $G$ . Therefore  $\phi$  fixes every element of  $G$  and hence  $\phi = 1$  (certainly we have  $\phi|_{S^2} = 1_{S^2}$ ). So  $G \triangleleft A$  as required.  $\square$

**Lemma 3.4** *Let  $S$  be a  $(2,5,5)$ -generating set of  $G$ . Then  $X = \text{Cay}(G, S)$  is a normal Cayley digraph.*

**Proof** Let  $S = \{a, c_1, c_2\}$  be a  $(2,5,5)$ -generating set of  $G$ , where  $a$  is an involution and  $c_1, c_2$  are elements of order 5.

If  $G = \langle a, c_i \rangle$ ,  $i = 1$  or  $i = 2$ , the lemma can be proved similarly to Lemma 3.3.

So we suppose that  $\langle a, c_i \rangle \neq G$  for  $i = 1, 2$ . Since  $G = \langle a, c_1, c_2 \rangle$ , we have  $\langle c_1 \rangle \neq \langle c_2 \rangle$ . Thus  $G = \langle c_1, c_2 \rangle$  and hence for the digraph  $\text{Cay}(G, \{c_1, c_2\})$ , the set  $\{c_1, c_2\}$  is *CI* by Theorem 1.4, and so  $\phi \in \text{Aut}(G)$  for  $\phi \in A_1$ . Therefore if  $\phi|_S = 1_S$ , then  $\phi|_{S^2} = 1_{S^2}$  and so (3.1) shows that  $X = \text{Cay}(G, S)$  is normal.  $\square$

**Lemma 3.5** *Let  $S$  be a  $(2,3,3)$ -generating set of  $G$ . Then  $X = \text{Cay}(G, S)$  is a normal Cayley digraph.*

**Proof** Assume that  $S = \{a, b_1, b_2 \mid a^2 = b_1^3 = b_2^3 = 1\}$ . If  $\langle a, b_i \rangle = G$  or  $\langle b_1, b_2 \rangle = G$ , for  $i = 1, 2$ , then the lemma can be proved similarly to Lemma 3.4. So it suffices to prove the result in the case:

- (i)  $\langle a, b_i \rangle \neq G$ ,  $i = 1, 2$ ; and
- (ii)  $\langle b_1, b_2 \rangle \neq G$ . (3.5)

Since  $A_5$  has only one conjugacy class of involutions, without loss of generality we may assume that  $a = (12)(34)$ . First we claim that if  $\langle a, b_1 \rangle \cong A_4$  then  $\langle a, b_2 \rangle \cong S_3$  by condition (3.5). Indeed, if  $\langle a, b_1 \rangle \cong A_4$  then we may suppose that  $b_1 = (123)$  and  $b_2 = (i, j, k)$ . Since  $G = \langle a, b_1, b_2 \rangle$ ,  $\{i, j, k\}$  must contain 5; say  $k = 5$ . Thus  $\{i, j\}$  cannot be one of  $\{1, 3\}$ ,  $\{2, 4\}$ ,  $\{2, 3\}$  or  $\{1, 4\}$  otherwise  $\langle a, b_2 \rangle \cong A_5$ . So  $\{i, j\}$  must be either  $\{1, 2\}$  or  $\{3, 4\}$ . In either case we have  $\langle a, b_2 \rangle \cong S_3$ . But if  $b_2 = (345)$  or  $(354)$  then  $\langle b_1, b_2 \rangle \cong A_5$ . From this it follows that  $b_2 = (125)$  or  $(152)$ . Similarly, suppose that  $a = (12)(34)$ ,  $b_2 = (125)$ . If  $\langle a, b_1 \rangle \cong S_3$ , then  $b_1$  must be  $(345)$  or  $(435)$ ; in either case we have  $\langle b_1, b_2 \rangle = G$  which contradicts condition (ii). Since  $\langle a, b_1 \rangle \neq G$ , it follows that  $\langle a, b_1 \rangle \cong A_4$ .

Finally it remains to show the result in the case that  $\langle a, b_1 \rangle \cong A_4$  and  $\langle a, b_2 \rangle \cong S_3$ , where  $a = (12)(34)$ ,  $b_1 = (123)$ , and  $b_2 = (125)$ .

Indeed, we claim that  $\phi|_S = 1_S$  for each  $\phi \in A_1$  in this case. If  $\phi|_S \neq 1_S$  then  $\phi$  fixes  $a$  and interchanges  $b_1$  and  $b_2$  and so  $\phi$  induces a graph automorphism between  $\text{Cay}(\langle a, b_1 \rangle, \{a, b_1\})$  and  $\text{Cay}(\langle a, b_2 \rangle, \{a, b_2\})$ . This contradicts the assumption  $\langle a, b_1 \rangle = A_4$  and  $\langle a, b_2 \rangle = S_3$ . Thus  $\phi|_S = 1_S$ . So  $\phi$  fixes  $a, b_1, b_2$  and hence  $A_1 = A_a = A_{b_1} = A_{b_2}$ . It follows that  $A_1 = 1$  and the lemma is proved.  $\square$

**Lemma 3.6** *Let  $S$  be a  $(2,2,3)$ -generating set of  $G$ . Then  $X = \text{Cay}(G, S)$  is a normal Cayley digraph.*

**Proof** In this case  $G$  is generated by a pair of involutions  $a_1, a_2$  and an element  $b$  of order 3. Without loss of generality we assume  $b = (123)$ . Let  $\phi \in A_1$ . If  $\phi|_S \neq 1_S$ , as before  $\phi$  fixes  $b$  and interchanges  $a_1$  and  $a_2$  and so it induces a graph automorphism

from  $\text{Cay}(\langle a_1, b \rangle, \{a_1, b\})$  to  $\text{Cay}(\langle a_2, b \rangle, \{a_2, b\})$ . So  $|\langle a_1, b \rangle| = |\langle a_2, b \rangle|$  and it follows that  $\langle a_1^\phi, b \rangle = \langle a_2, b \rangle$ .

Now we consider three cases.

(1)  $\langle a_1, b \rangle \cong \langle a_2, b \rangle \cong A_5$ .

Since  $\text{Cay}(\langle a_i, b \rangle, S_i)$  is a normal Cayley digraph by Theorem 1.4 (where  $S_i = \{a_i, b\}$  and  $i = 1, 2$ ), we have  $\phi \in \text{Aut}(G, \{a_i, b\}) \leq \text{Aut}(G)$  and so the result is proved by Lemma 1.3.

(2)  $\langle a_1, b \rangle \cong \langle a_2, b \rangle \cong A_4$ .

It is obvious that  $\langle a_1, b \rangle \neq \langle a_2, b \rangle$ . Let  $\phi \in A_1$ . If  $\phi|_S = 1_S$  then  $\phi$  induces an automorphism of the subgraph  $X_i = \text{Cay}(\langle a_i, b \rangle, S_i)$ , where  $S_i = \{a_i, b\}$  and  $i = 1, 2$ . Since  $C_1 = (1, a_1)$  is the unique circuit of length 2 and  $C_2 = (1, b, b^2)$  is the unique circuit of length 3 at the point 1, it follows that  $\phi$  fixes  $C_1$  and  $C_2$  pointwise. Since  $\phi$  fixes  $b$ ,  $\phi$  fixes the neighbourhood of  $b$  in  $X_1$  and thus  $\phi$  fixes  $a_1b$ . Since  $\phi$  fixes  $a_1$ ,  $\phi$  fixes the neighbourhood of  $a_1$  in  $X_1$  and thus  $\phi$  fixes  $ba_1$ . Since  $\phi$  fixes  $a_1$ ,  $\phi$  fixes the neighbourhood of  $a_1$  in  $X$  and thus  $\phi$  fixes  $a_2a_1$ . Similarly we can check that  $\phi$  fixes the points  $a_2b, a_2^2, ba_2$  and  $a_1a_2$ . This shows that  $\phi|_{S_2} = 1_{S_2}$ . By Lemma 3.1,  $\text{Cay}(G, S)$  is normal as required.

(3)  $\langle a_1, b \rangle \cong \langle a_2, b \rangle \cong S_3$ .

This case can not happen, since for each element of order 3 in the unique subgroup which is isomorphic to  $S_3$ , if  $\langle a_1, b \rangle \cong \langle a_2, b \rangle \cong S_3$  then  $\langle a_1, b \rangle = \langle a_2, b \rangle \cong S_3$  which contradicts  $\langle a_1, a_2, b \rangle = G$ .

By (1), (2) and (3), the lemma is proved.  $\square$

**Lemma 3.7** *Let  $S$  be a  $(2, 2, 5)$ -generating set of  $G$ . Then  $X = \text{Cay}(G, S)$  is a normal Cayley digraph.*

**Proof** Suppose that  $G$  is generated by a pair of involutions  $a_1, a_2$  and an element  $c$  of order 5. If  $\langle a_1, c \rangle \cong D_{10}$  then  $\langle a_2, c \rangle \not\cong D_{10}$  by the fact that  $c$  is in a unique subgroup  $H \cong D_{10}$ , and so  $\langle a_2, c \rangle \cong A_5$ . Now we distinguish two cases.

(1)  $\langle a_1, c \rangle \cong D_{10}$  and  $\langle a_2, c \rangle \cong A_5$ ;

(2)  $\langle a_1, c \rangle \cong \langle a_2, c \rangle \cong A_5$ .

The lemma follows in either case in the same way as the proof of Lemma 3.5.  $\square$

**Lemma 3.8** *Let  $S$  be a  $(5, 5, 5)$ -generating set of  $G$ . Then  $X = \text{Cay}(G, S)$  is a normal Cayley digraph.*

**Proof** Let  $S = \{c_1, c_2, c_3 \mid c_1^5 = c_2^5 = c_3^5 = 1\}$ . Clearly if  $\langle c_1 \rangle \neq \langle c_2 \rangle$ , then  $\langle c_1, c_2 \rangle = A_5$ . Since  $G = \langle c_1, c_2, c_3 \rangle$ , we may assume that  $\langle c_1, c_2 \rangle = A_5$ . If  $\langle c_2, c_3 \rangle \cong Z_5$ , then  $\phi|_S = (c_1, c_2, c_3)$  is not a graph automorphism from  $\text{Cay}(G, \{c_1, c_2\})$  to  $\text{Cay}(G^\phi, \{c_1, c_2\}^\phi)$  since  $\langle c_1, c_2 \rangle \cong A_5$  and  $\langle c_1, c_2 \rangle^\phi = \langle c_2, c_3 \rangle \cong Z_5$ .

So we assume that  $\langle c_1, c_2 \rangle = \langle c_2, c_3 \rangle = \langle c_3, c_1 \rangle = G$ . Let  $A_{1, c_1}$  denote the subgroup of  $A$  which fixes 1 and  $c_1$ . Then  $|A_1| \leq 3|A_{1, c_1}|$ . Let  $\phi \in A_{1, c_1}$ , then  $\phi$  fixes  $c_1$  and stabilizes the set  $\{c_2, c_3\}$ . Set  $S_1 = \{c_2, c_3\}$ . Thus  $\phi$  induces a graph automorphism of  $\text{Cay}(G, S_1)$ . Since  $\text{Cay}(G, S_1)$  is a normal Cayley graph by Theorem 1.4, it follows that  $\phi \in \text{Aut}(G)$ . But  $\text{Aut}(G)$  has no non identity automorphism that interchanges

a pair of elements  $(c_2$  and  $c_3)$  of order 5 and fixes another element  $c_1$  of order 5, so we have  $\phi = 1$  and thus  $A_{1,c_1} = 1$  and  $|A_1| \leq 3$ .

If  $|A_1| = 3$ , then  $|A| = 180$  and  $A$  is not simple group. So  $A$  contains a non-trivial normal subgroup  $N$ , that is  $N \neq 1$  and  $N \neq A$ . If  $G \cap N \neq 1$  then  $G \cap N$  is a nontrivial normal subgroup of  $G$  which contradicts  $G \cong A_5$ . It follows that  $G \cap N = 1$  and thus  $A = N \cdot G$  and  $A_1 = N$ . Since  $A_1$  is the stabilizer of the point 1, we get  $N = 1$ , a contradiction. It follows that  $|A_1| \leq 2$  and  $\text{Cay}(G, S)$  is a normal Cayley digraph as claimed.  $\square$

**Lemma 3.9** ([4, Lemma 2.2]) *Let  $X = \text{Cay}(G, S)$ . Then  $X$  is a normal Cayley digraph of  $G$  if the following conditions hold:*

- (i) *for each  $\phi \in A_1$  there exists  $\sigma \in \text{Aut}(G)$  such that  $\phi|_S = \sigma|_S$ ;*
- (ii) *for each  $\phi \in A_1$ ,  $\phi|_S = 1_S$  implies  $\phi|_{S^2} = 1_{S^2}$ .*

**Proof** (1) Condition (ii) implies that if  $\phi \in A_1$  and  $\phi|_S = 1_S$ , then  $\phi = 1_G$ .

(2) We show that  $A_1 \leq \text{Aut}(G, S)$ . By the hypothesis (i), for each  $\phi \in A_1$ , we may take  $\sigma \in \text{Aut}(G)$  such that  $\phi|_S = \sigma|_S$ . Then  $\phi\sigma^{-1}|_S = 1_S$ . By the proof above we have  $\phi\sigma^{-1} = 1_G$  and  $\phi = \sigma \in \text{Aut}(G, S)$ .

(1) and (2) imply that  $X$  is a normal graph of the group  $G$ .  $\square$

**Lemma 3.10** *Let  $S$  be a  $(3,3,3)$ -generating set of  $G$ . Then  $X = \text{Cay}(G, S)$  is a normal Cayley digraph.*

**Proof** First suppose that  $G$  is generated by three elements  $b_1, b_2$  and  $b_3$  of order 3. We now distinguish two cases:

(i) There exist  $b_i, b_j \in S$  such that  $\langle b_i, b_j \rangle = G$ . In this case there must exist  $b_i$  or  $b_j$  such that  $\langle b_i, b_k \rangle \neq G$  or  $\langle b_j, b_k \rangle \neq G$ . As in Lemma 3.8 we can prove  $\text{Cay}(G, S)$  is a normal digraph of  $G$ .

(ii) There are no  $b_i, b_j \in S$  such that  $\langle b_i, b_j \rangle = G$ . In this case there exist  $b_i, b_j$  such that  $\langle b_i, b_j \rangle = A_4$ . It is clear that the 3-cycle  $b_i$  and  $b_j$  are in a same subgroup of  $G$  which is isomorphic to  $A_4$  if and only if they have two symbols that are the same. So without loss of generality, we assume that  $b_1 = (123)$  and  $b_2 = (124)$ . Since  $\langle b_1, b_2, b_3 \rangle = A_5$ , thus  $b_3 = (i, j, 5)$ , and in addition  $\langle b_i, b_3 \rangle \cong A_4$ , for  $i = 1, 2$  (If  $\langle b_i, b_3 \rangle \cong Z_3$ , then  $b_3 \in \langle b_i, b_j \rangle \cong A_4$ , a contradiction.) Hence we have  $b_3 = (125)$  or  $(152)$ .

In this case for each  $\phi \in A_1$ , if  $\phi|_S = (b_1, b_2, b_3)$ , then there exists  $\sigma \in \text{Aut}(G, S)$  such that  $\phi|_S = \sigma|_S$ ; in fact by the assumption,  $\sigma = (345)$ . If  $\phi|_S = 1_S$ , then  $\phi$  induces a graph automorphism of  $\text{Cay}(\langle b_i, b_j \rangle, \{b_i, b_j\})$  for  $i, j \in \{1, 2, 3\}$ . Set  $S_1 = \{b_i, b_j\}$ . It is easy to check that the Cayley graph  $\text{Cay}(\langle b_i, b_j \rangle, S_1)$  satisfies the conditions of Lemma 3.9. So  $\text{Cay}(\langle b_i, b_j \rangle, S_1)$  is normal and hence  $\phi|_{S_1} \in \text{Aut}(\langle b_i, b_j \rangle)$ . It follows that  $\phi$  fixes  $b_i$  and  $b_j$  and hence fixes  $b_i^2, b_j^2, b_i b_j$  and  $b_j b_i$ . Similarly  $\phi$  fixes the other elements of  $S^2$ . So if  $\phi|_S = 1_S$  then  $\phi|_{S^2} = 1_{S^2}$ . By Lemma 3.9 our statement follows.

By (i) and (ii), the lemma is proved.  $\square$



Now consider the case when  $G$  is generated by three involutions  $a_1, a_2$  and  $a_3$  and  $X = \text{Cay}(G, S)$  is an undirected graph. We will use the following result.

**Lemma 3.11** ([5, Theorem 1.3]) *Suppose that  $G$  is a nonabelian simple group. Then  $G$  is a 3-CI-group if and only if  $G = A_5$ .*

**Lemma 3.12** *Let  $S$  be a  $(2,2,2)$ -generating set of  $G$ . Then  $X = \text{Cay}(G, S)$  is a normal Cayley graph.*

**Proof** Since  $S = \{a_1, a_2, a_3\}$  is 3-CI, for each  $\phi \in A_1$  we have  $\phi \in \text{Aut}(G)$  by Lemma 3.11 and the lemma follows.  $\square$

**PROOF OF THEOREM 1.5** This is given by Lemmas 3.2, 3.3, 3.4, 3.5, 3.6, 3.7, 3.8, 3.10 and 3.12.  $\square$

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